Programming Models

Types of parallelism

- Data Parallelism
 - Each processor performs the same task on different data (remember SIMD, MIMD)
- Task Parallelism
 - Each processor performs a different task on the same data (remember MISD, MIMD)
- Many applications incorporate both

Implementation: Single Program Multiple Data

- Dominant programming model for shared and distributed memory machines
- One source code is written
- Code can have conditional execution based on which processor is executing the copy
- All copies of code start simultaneously and communicate and synchronize with each other periodically



Data Parallel Programming

- Example
 - One code will run on 2 CPUs
 - Program has array of data to be operated on by 2 CPUs so array is split into two parts.

program: 	CPU A	CPU B
<pre>if CPU=a then low_limit=1 upper_limit=50 elseif CPU=b then low_limit=51 upper_limit=100 end if do I = low_limit, upper_limit work on A(I) end do end program</pre>	<pre>program: low_limit=1 upper_limit=50 do I= low_limit, upper_limit work on A(I) end do end program</pre>	<pre>program: low_limit=51 upper_limit=100 do I= low_limit, upper_limit work on A(I) end do end program</pre>

Task Parallel Programming

- Example
 - One code will run on 2 CPUs
 - Program has 2 tasks (a and b) to be done by 2 CPUs



Shared Memory Programming: pthreads

- Shared memory systems (SMPs, ccNUMAs) have a single address space
- Applications can be developed in which loop iterations (with no dependencies) are executed by different processors
- Threads are 'lightweight processes' (same PID)
- Allows 'MIMD' codes to execute in shared address space

Shared Memory Programming: OpenMP

- Built on top of pthreads
- Shared memory codes are mostly data parallel, 'SIMD' kinds of codes
- OpenMP is a standard for shared memory programming (compiler directives)
- Vendors offer native compiler directives

Accessing Shared Variables

- If multiple processors want to write to a shared variable at the same time, there could be conflicts :
 - Process 1 and 2
 - read X
 - compute X+1
 - write X



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 Programmer, language, and/or architecture must provide ways of resolving conflicts (mutexes and semaphores)

OpenMP Example #1: Parallel Loop

!\$OMP PARALLEL DO
do i=1,128
 b(i) = a(i) + c(i)
end do
!\$OMP END PARALLEL DO

```
void simple(int n, float *a, float *b)
{
    int i;
#pragma omp parallel for
    for (i=1; i<n; i++) /* i is private by default */
        b[i] = (a[i] + a[i-1]) / 2.0;</pre>
```

- The first directive specifies that the loop immediately following should be executed in parallel.
- The second directive specifies the end of the parallel section (optional).
- For codes that spend the majority of their time executing the content of simple loops, the PARALLEL DO directive can result in significant parallel performance.

OpenMP Example #2: Private Variables

```
!$OMP PARALLEL DO SHARED(A,B,C,N) PRIVATE(I,TEMP)
do I=1,N
    TEMP = A(I)/B(I)
    C(I) = TEMP + SQRT(TEMP)
end do
!$OMP END PARALLEL DO
```

- In this loop, each processor needs its own private copy of the variable TEMP.
- If TEMP were shared, the result would be unpredictable since multiple processors would be writing to the same memory location.



Distributed Memory Programming: MPI

- Distributed memory systems have separate address spaces for each processor
- Local memory access is faster than remote memory
- Data must be manually decomposed
- MPI is the de facto standard for distributed memory programming (library of subprogram calls)
- Vendors typically have native libraries such as SHMEM (T3E) and LAPI (IBM)



Data Decomposition

- For distributed memory systems, the 'whole' grid is decomposed to the individual nodes
 - Each node works on its section of the problem

Grid of Problem to be solved



Typical Data Decomposition

• Example: integrate 2-D propagation problem:



MPI Example #1

```
    Every MPI program needs these:

   #include "mpi.h"
   int main(int argc, char *argv[])
   int nPEs, iam;
   /* Initialize MPI */
   ierr = MPI Init(&argc, &argv);
   /* How many total PEs are there */
   ierr = MPI Comm size(MPI COMM WORLD, &nPEs);
   /* What node am I (what is my rank?) */
   ierr = MPI Comm rank(MPI COMM WORLD, &iam);
   ...
   ierr = MPI Finalize();
```

MPI Example #2

```
#include "mpi.h"
int main(int argc, char *argv[])
int numprocs, myid;
MPI Init(&argc,&argv);
MPI Comm size(MPI COMM WORLD,&numprocs);
MPI Comm_rank(MPI_COMM_WORLD,&myid);
/* print out my rank and this run's PE size */
printf("Hello from %d of %d\n", myid, numprocs);
MPI Finalize();
```

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MPI: Sends and Receives

- MPI programs must send and receive data between the processors (communication)
- The most basic calls in MPI (besides the three initialization and one finalization calls) are:
 - MPI_Send
 - MPI_Recv
- These calls are blocking: the source processor issuing the send/receive cannot move to the next statement until the target processor issues the matching receive/send.

Message Passing Communication

- Processes in message passing programs communicate by passing messages
- Basic message passing primitives: MPI_CHAR, MPI_SHORT, ...

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- Send (parameters list)
- Receive (parameter list)
- Parameters depend on the library used
- Barriers

MPI Example #3: Send/Receive

```
#include "mpi.h"
int main(int argc, char *argv[])
    int numprocs, myid, tag, source, destination, count, buffer;
    MPI Status status;
    MPI Init(&argc,&argv);
    MPI Comm size (MPI COMM WORLD, & numprocs);
    MPI Comm rank (MPI COMM WORLD, &myid);
    tag=1234;
    source=0;
    destination=1;
    count=1;
    if(myid == source) {
        buffer=5678;
        MPI Send(&buffer,count,MPI INT,destination,tag,MPI COMM WORLD);
        printf("processor %d sent \sqrt[8]{d}n", myid, buffer);
    if(myid == destination) {
        MPI Recv(&buffer,count,MPI INT,source,tag,MPI COMM WORLD,&status);
        printf("processor %d got %d\n",myid,buffer);
    MPI Finalize();
ł
```

QUESTIONS?

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