



High speed computer networks

Frame relay

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Introduction

□ Switching

- ✓ Circuit switching
- ✓ Packet switching
 - Virtual circuit
 - Datagram

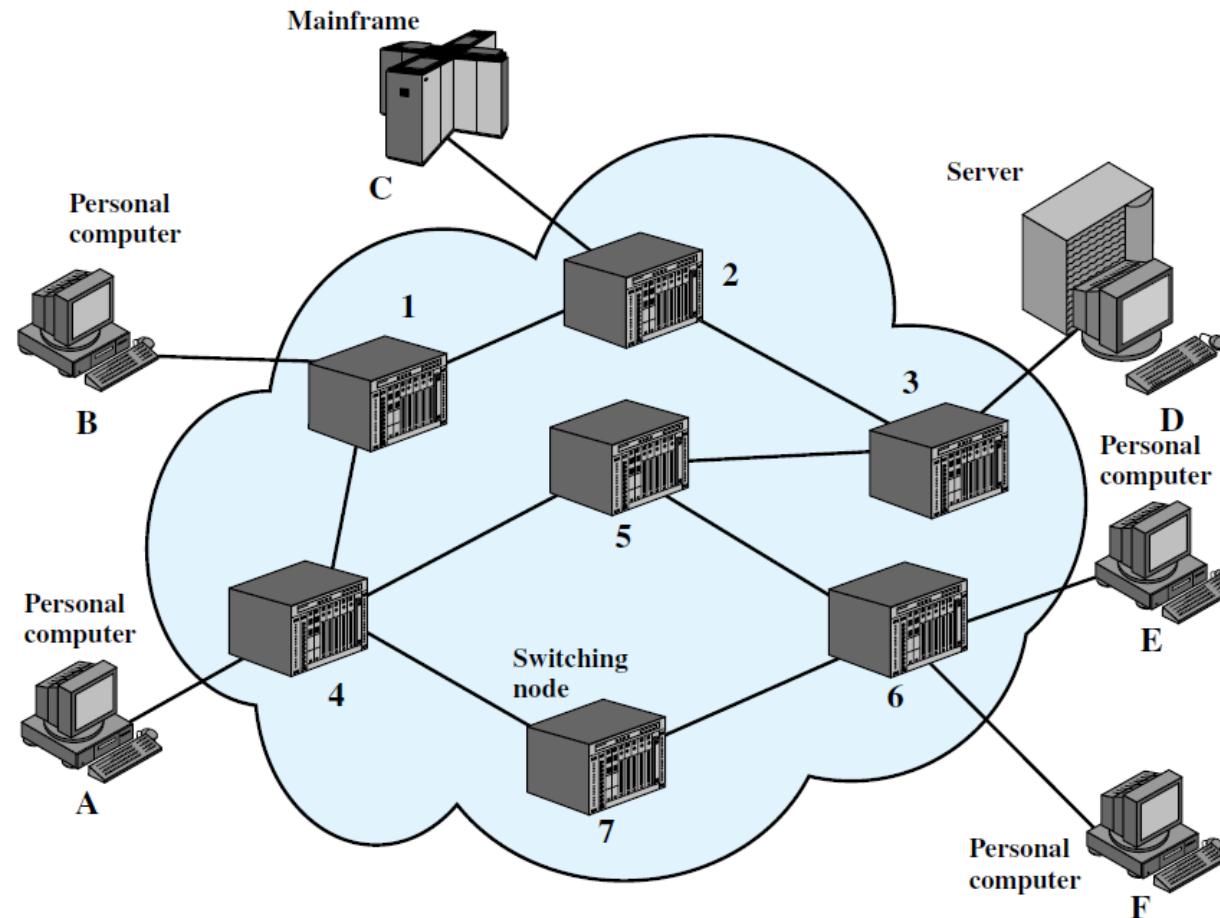
□ X.25

□ Frame relay networks

- ✓ Architecture

Switched communication network

- Transmitting data through a network of intermediate switching nodes



Circuit-Switching

- ❑ Designed to handle voice traffic

- ❑ Process

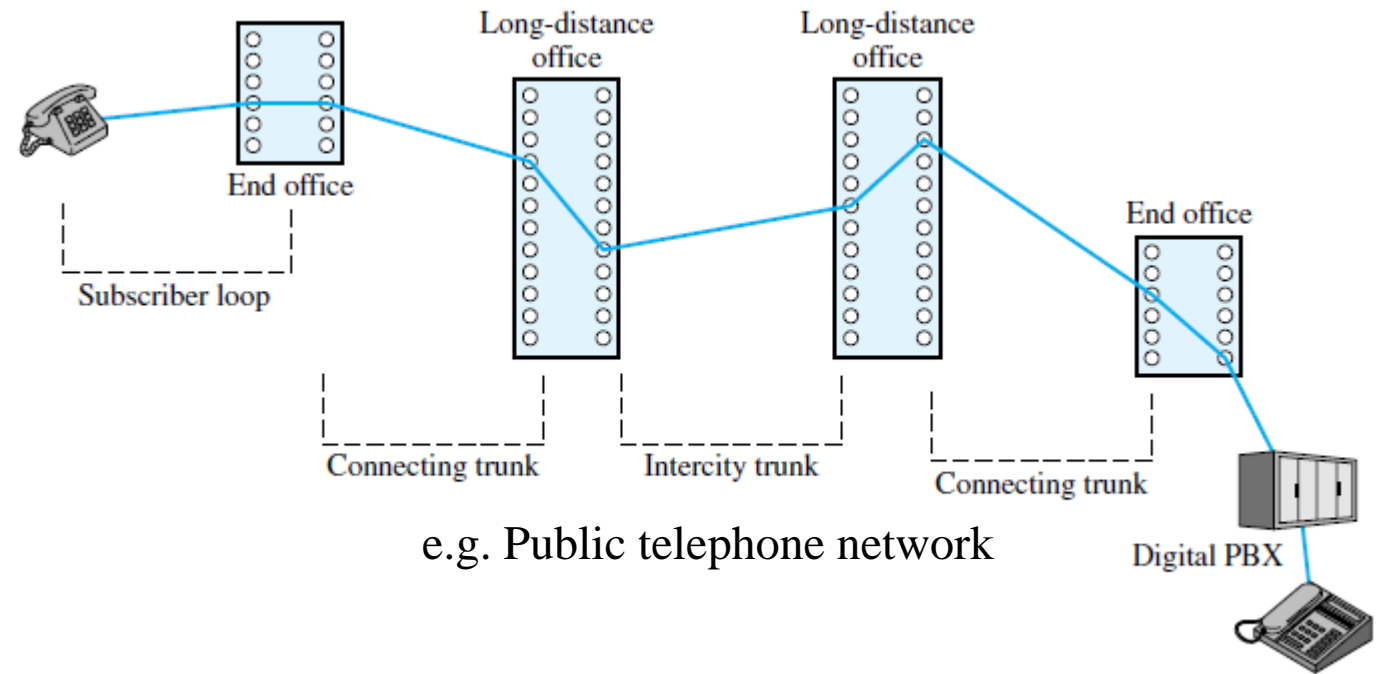
- ✓ End-to-end circuit establishment
- ✓ Data transfer
- ✓ Circuit disconnect

- ❑ Advantage

- ✓ No variation of delay

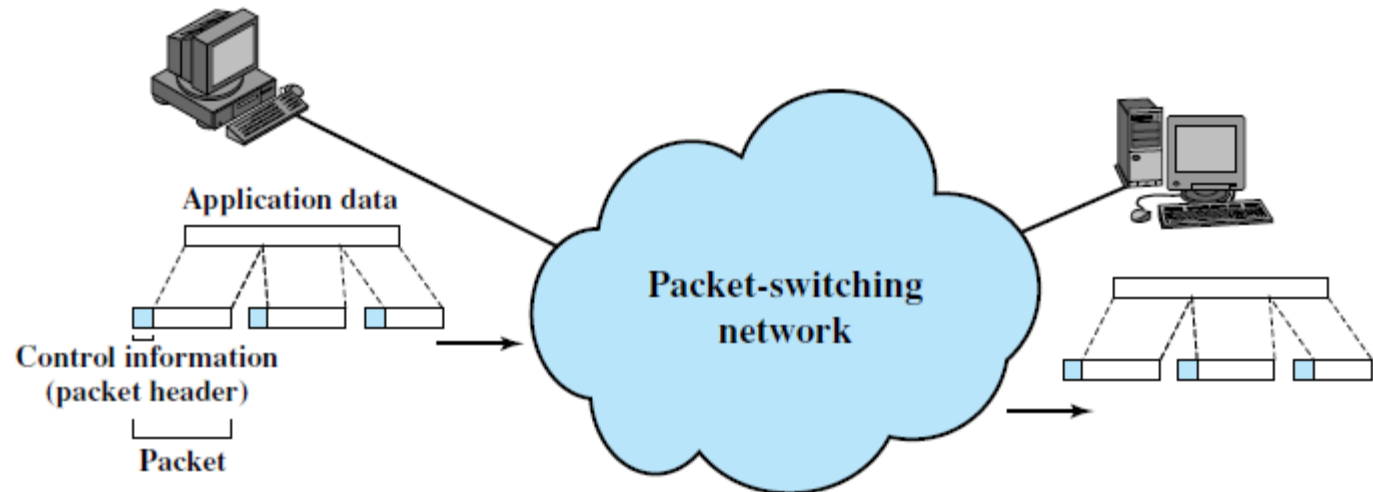
- ❑ Shortcomings

- ✓ Inefficient
 - high idle time
- ✓ Constant data rate
 - devices must transmit and receive at the same data rate



Packet switching

- ❑ Used extensively for computer communications
- ❑ Data transmitted in short blocks, or packets
 - ✓ Packet length < 1000 octets
 - ✓ Each packet contains user data plus control info (routing)
 - ✓ Store and forward



Packet switching

□ Path selection- considering the state of the network

□ Advantages

- ✓ Link can be shared dynamically by many packets
- ✓ Different data rate can be used
 - A packet-switching network can perform data-rate conversion
- ✓ When congestion occurs, packets can be still accepted
 - But, delay time increases
- ✓ Priority can be used

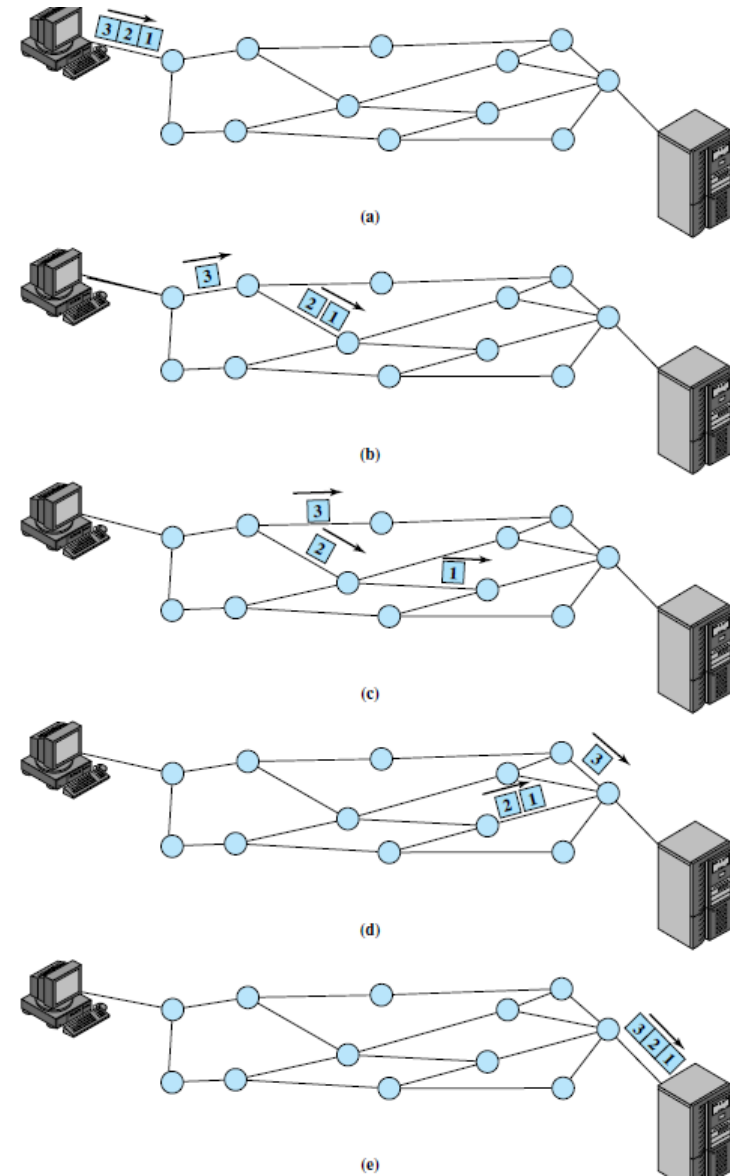
Packet switching

❑ Disadvantages

- ✓ Packets incur additional delay with every node they pass through
- ✓ Jitter: variation in packet delay
- ✓ Data overhead in every packet for routing information
- ✓ Processing overhead for every packet at every node traversed

Packet switching : switching Technique

- ❑ How the network will handle this stream of packets ?
- ❑ **Datagram approach (connectionless)**
 - ✓ Each packet is treated independently
 - ✓ The packets do not all follow the same route
 - ✓ The exit node or the destination node restores the packets to their original order
- ❑ Each packet must contain a destination address and possibly a source address
 - ✓ Efficient for fairly long packets

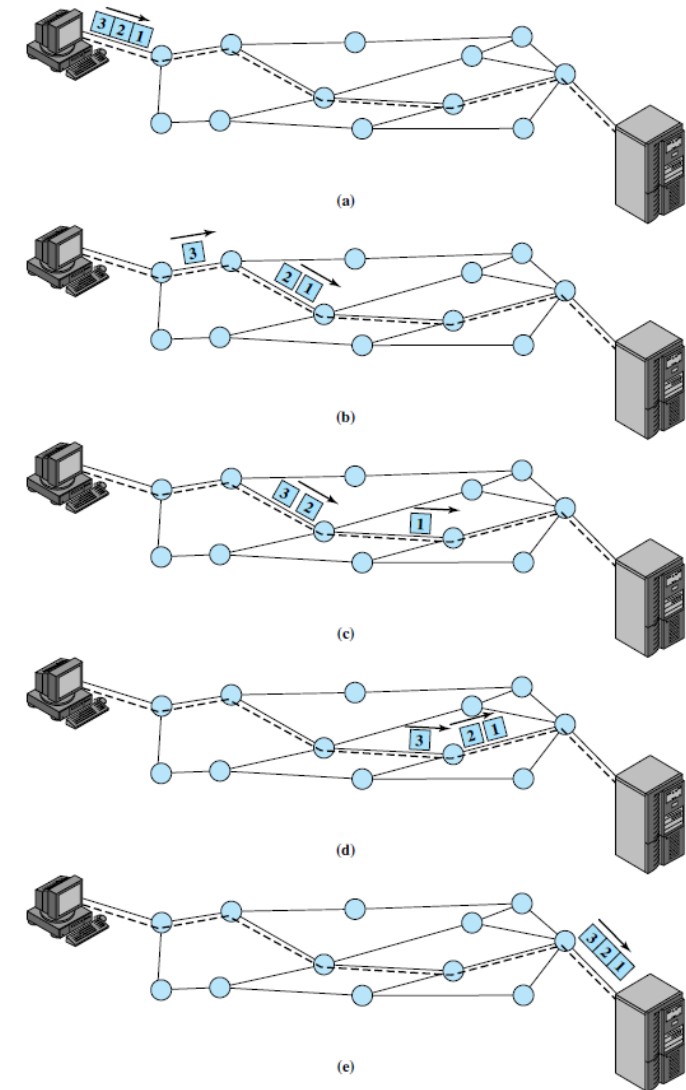


Virtual circuit approach

❑ Connection-oriented

- ✓ Pre-defined route is established
- ✓ Exchange of control packets prior to the transfer of traffic
- ✓ All packets follow the same route
- ✓ Each packet contains data plus a virtual circuit identifier
- ✓ No routing decision

❑ Any difference with circuit switching?



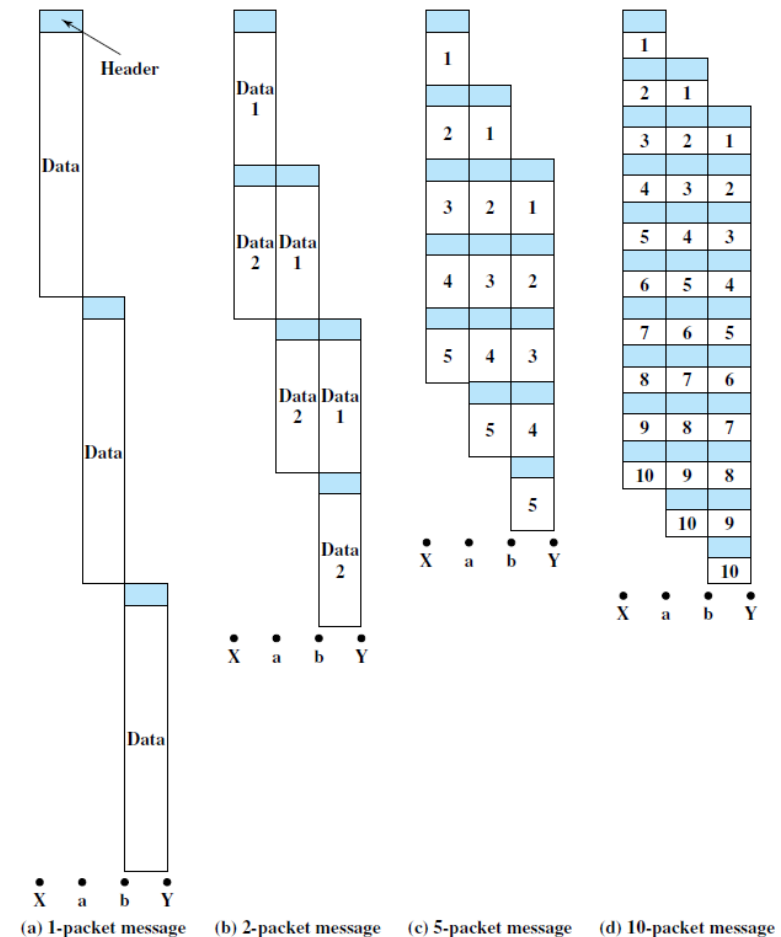
Virtual circuit approach

- ❑ If two stations wish to exchange **data over an extended period of time**, there are certain advantages to virtual circuits
 - ✓ The network may provide services such as sequencing and error control
 - ✓ Fast packet transfer

- ❑ But, datagram is preferred for exchanging only a few packets
 - ✓ No call set up phase
 - ✓ Also flexible and reliable

Packet size and transmission time

- ❑ Relation between packet size and transmission time
- ❑ E.g., assume a virtual circuit from station X to Y
 - ✓ Message to be sent 40 byte and 3 bytes control information
 - ✓ Break the message into two packet blocks
 - ✓ Continue decreasing the block size ?



X.25

- ❑ A standard (ITU-T) used for traditional packet switching networks
- ❑ Specifies an interface between a host system and a packet switching network
 - ✓ Set of protocols corresponding to the first three layers of OSI

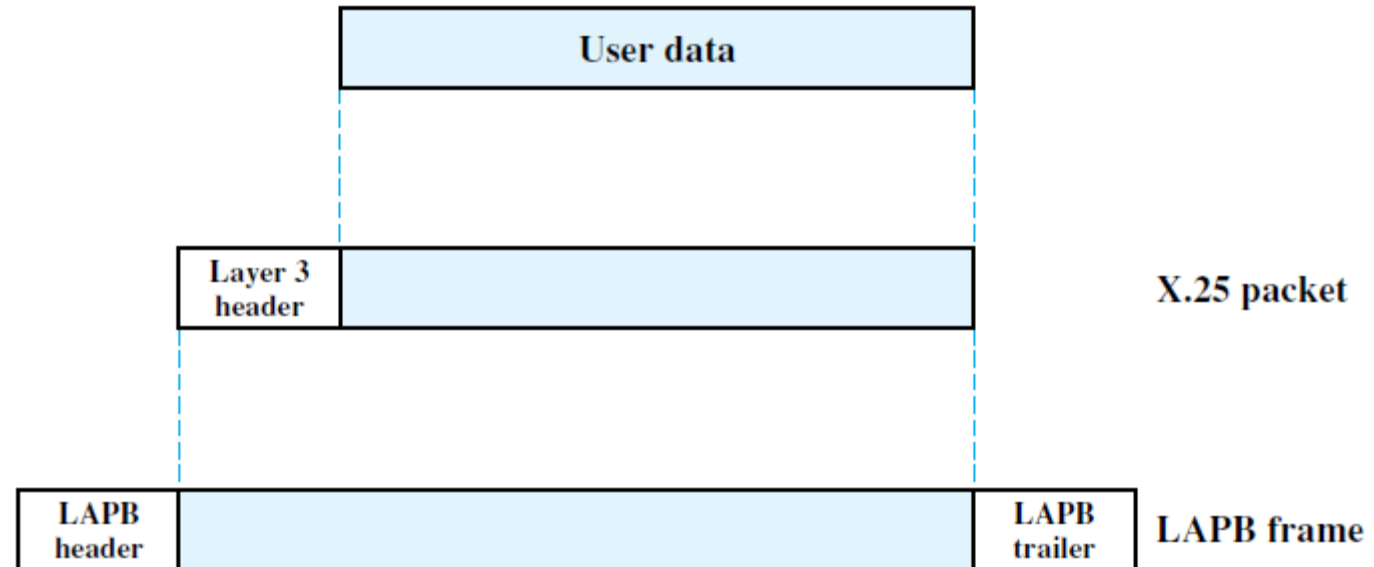
- ❑ Physical layer
 - ✓ Interface between stations and packet-switching network
 - ✓ Use Physical layer specifications in X.25

- ❑ Link layer
 - ✓ Reliable transmission of data between packet switches (link access protocol-balanced LAP-B)

- ❑ Packet layer
 - ✓ Packet formats, control procedures to setup a call and exchange information
 - ✓ logical connection between two stations through the network

X.25

□ X.25 protocol control information



Frame relay networks: high speed networking

- ❑ Designed to provide a more efficient transmission scheme than X.25

- ❑ Key features of X.25
 - ✓ Call control packets are carried on the same channel and virtual circuit as data packet
 - ✓ Multiplexing of virtual circuits takes place at layer 3
 - ✓ Both layer 2 and layer 3 include flow control and error control mechanisms
 - Large overhead and slow down transmission

Great improvement on the reliability of physical circuits (e.g. bit error rate of a typical optical link is 10^{-8})

➤ **Remove unnecessary overhead**

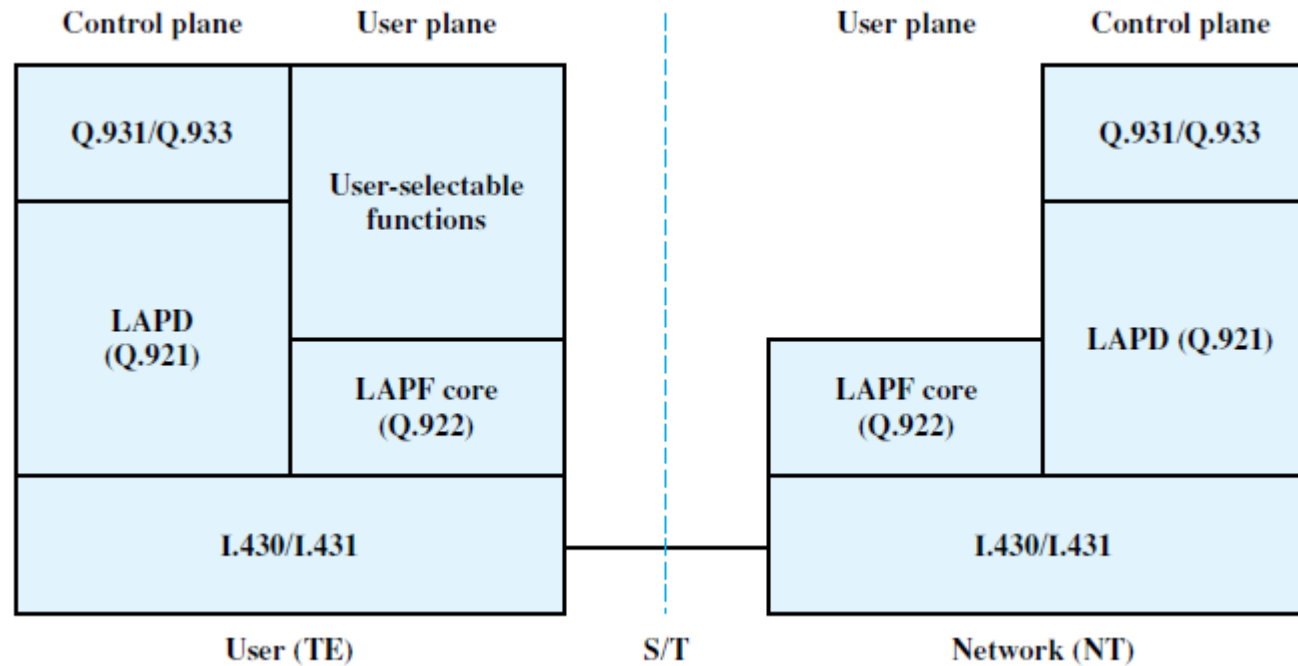
Frame relay

- ❑ Removes overhead that X.25 imposes and takes advantage of the improvement in transmission link

- ❑ Key differences
 - ✓ Call control signaling on separate logical connection from user data
 - ✓ Multiplexing/switching of logical connections at layer 2 (not layer 3)
 - ✓ No hop-by-hop flow control and error control
 - Flow and error control are the responsibility of higher layers
 - ✓ Throughput an order of magnitude higher than X.25

Protocol architecture

- Frame Relay has 2 layers: physical and data link (or LAPF)



Protocol architecture

□ Control plane

- ✓ Signaling over D channel
- ✓ Separate channel for control information
- ✓ At the data link layer LAPD (link access control –D channel) is used –for reliable data link control

□ User plane

- ✓ For actual transfer of information
- ✓ LAPF(link access procedure for frame mode bearer service)

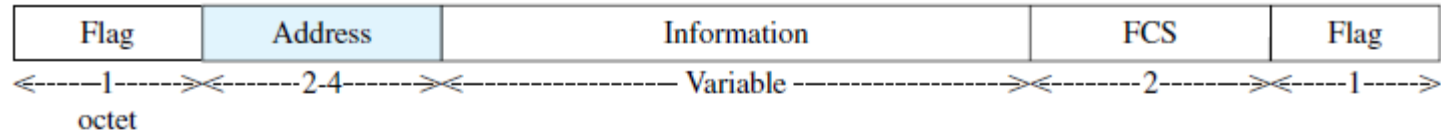
- Detection of transmission errors
- Virtual circuit multiplexing/demultiplexing
- Frame delimiting, alignment and transparency
- Congestion control function



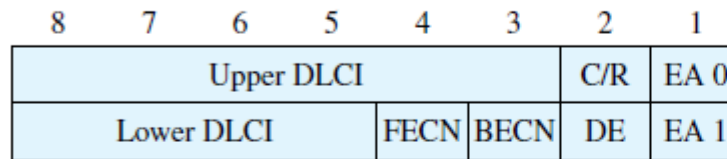
Also functions
of LAPD

User data transfer

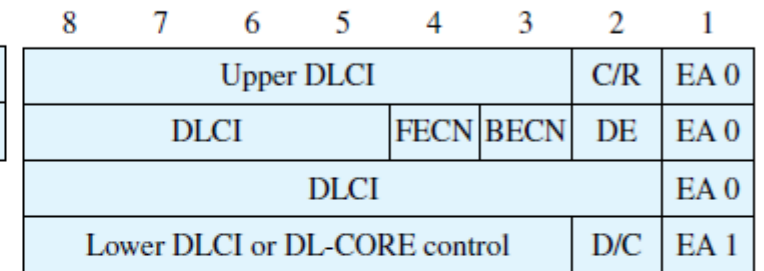
- LAPF core protocol
 - ✓ Similar to LAPD



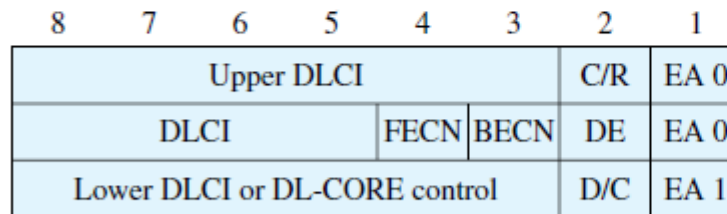
(a) Frame format



(b) Address field—2 octets (default)



(d) Address field—4 octets



(c) Address field—3 octets

- EA Address field extension bit
- C/R Command/response bit
- FECN Forward explicit congestion notification
- BECN Backward explicit congestion notification
- DLCI Data link connection identifier
- D/C DLCI or DL-CORE control indicator
- DE Discard eligibility

User data transfer

- ❑ No control field, which is normally used for
 - ✓ Identify frame type (data or control)
 - ✓ Sequence numbers
- ❑ Implication:
 - ✓ Connection setup/teardown carried on separate channel
 - ✓ Cannot do flow and error control

- ❑ Flags – mark the beginning and end of frame
- ❑ Address – 2 bytes (can be extended to 3 or 4 bytes)
 - ✓ DLCI – same as the virtual circuit number
 - ✓ Forward explicit congestion notification (FECN) and backward explicit congestion notification (BECN) – used as part of frame relay congestion control process

Congestion control

❑ Discard control (DE)

- ✓ When set, indicates that a frame should be discarded in preference to other
- ✓ Part of traffic management procedures
 - Committed burst size – the maximum amount of data that a network agrees to transfer under normal circumstance
 - Based on Committed information rate (CIR) – in the event of CIR being exceeded
 - Take no action, set the DE field, discard the frame immediately

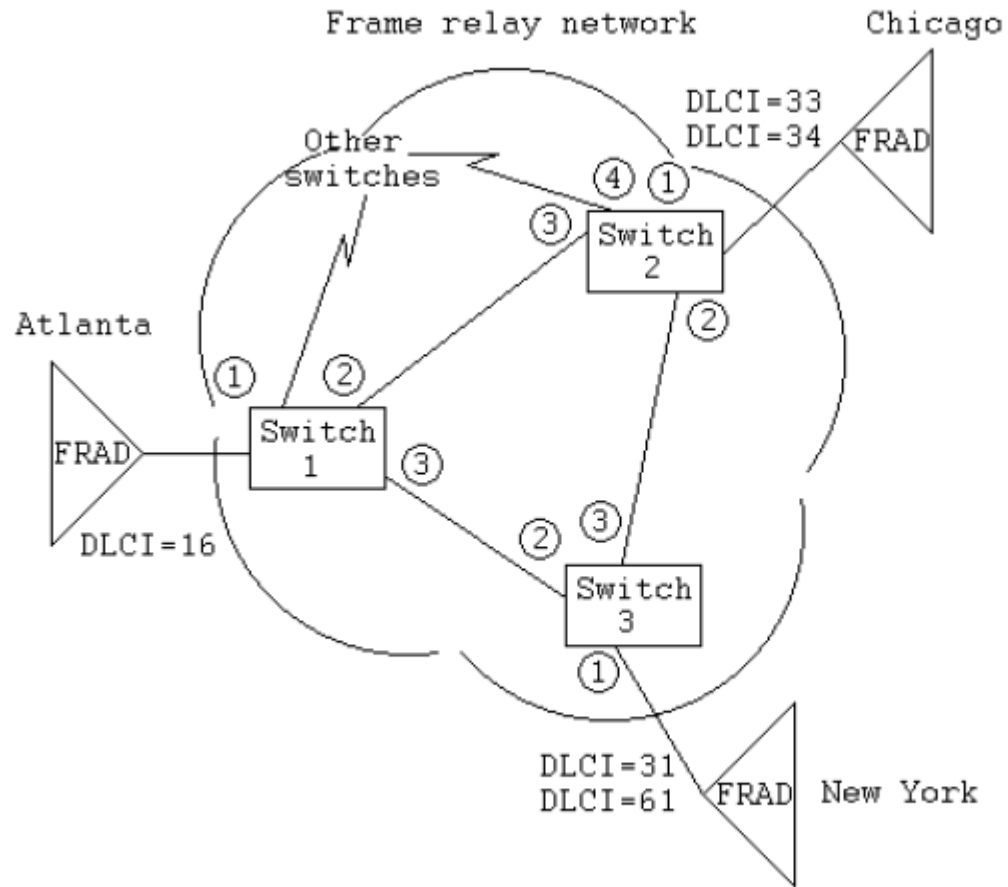
❑ BECN

- ✓ Informs the sender that congestion has occurred
- ✓ The source - slow down to prevent the loss of packets

❑ FECN

- ✓ Congestion control in the direction of the data transmission
- ✓ The forward explicit congestion notification (FECN) bit can be set by any switch to indicate that traffic is congested
- ✓ The destination knows that it should expect delay or a loss of packets

DLCI



Input		Output	
Port	DLCI	Port	DLCI
1	16	3	xyz
3	abc	1	16

Frame relay call control

□ Data transfer involves:

- ✓ Establish logical connection and DLCI
- ✓ Exchange data frames
- ✓ Release logical connection

□ 4 message types needed

- ✓ SETUP
- ✓ CONNECT
- ✓ RELEASE
- ✓ RELEASE COMPLETE

Virtual circuit

❑ Permanent virtual circuit

- ✓ Permanently established connection
- ✓ No call setup and termination
- ✓ For frequent and consistent data transfers

❑ Switched virtual circuit

- ✓ Creates temporary short connection
- ✓ Establishing by sending signaling messages to the network